

# AVal: an Extensible Attribute-Oriented Programming Validator for Java

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## Abstract

*Attribute Oriented Programming (@OP) permits programmers to extend the semantics of a base program by annotating it with attributes that are related to a set of concerns. Examples of this are applications that rely on XDoclet (such as Hibernate) or, with the release of Java5's annotations, EJB3. The set of attributes that implements a concern defines a Domain Specific Language, and as such, imposes syntactic and semantic rules on the way attributes are included in the program or even on the program itself. We propose a framework for the definition and checking of these rules for @OP that uses Java5 annotations. We define an extensible set of meta-annotations to allow the validation of @OP programs, as well as the means to extend them using a compile-time model of the program's source code. We show the usefulness of the approach by presenting two examples of its use: an @OP extension for the Fractal component model called Fraclet, and the JSR 181 for web services definition.*

## 1. Introduction

Attribute Oriented Programming[14] (@OP) allows the extension of a base language's semantics with concepts specific to a given domain by means of attributing the source code elements of the language with relevant metadata. For example, a class with an attribute `Persistent` may specify that the instances of that class must be saved in a persistent storage at runtime. Attribute-Oriented Programming can be implemented through library support (XDoclet, Apache Commons Attributes), or by direct language support (.Net, Java).

A set of attributes in an @OP framework can be seen as defining a Domain Specific Language (@DSL) on top of the base language. This @DSL imposes a set of syntactic and semantic rules in addition to those of the base language. The validations required to check these rules vary among @OP

frameworks, and they can be quite complex. In the case of JSR-181<sup>1</sup>, some annotations reference XML descriptors that must be validated by the framework; while for other annotations, validations must be performed on the code elements on which the annotations are placed; for example, a `@WebService` object cannot implement the `finalize()` method.

Although some support is included for the validation of these rules in existing frameworks and languages, in general it is not enough to cover the complex rules specified in current @OP languages. It is then up to the framework programmer to implement these supplementary checks in a way that is expressive, extensible, and provides meaningful error messages to the end user. In this paper, we present a framework for the validation of @OP programs which covers these properties.

### 1.1. Attribute Oriented Programming

Attribute-oriented programming is a program-level marking technique. Basically, developers can mark program elements (*e.g.*, classes, methods and fields) with attributes (*annotations*) to indicate that they maintain application-specific or domain-specific semantics [14]. Annotations separate the application's business logic from middleware-specific or domain-specific semantics (*e.g.*, logging and web service functions).

By hiding the implementation details of those semantics from program code, annotations increase the level of programming abstraction and reduce programming complexity, resulting in simpler and more readable programs. The program elements associated with annotations are transformed to more detailed programs by a supporting tool (*e.g.*, generation engine). For example, a generation engine may insert a call to a logging API into the methods associated with a `logging` annotation. The dependencies towards the underlying API are replaced by annotations, acting as weak

<sup>1</sup>Java Specification Request 181 for web service metadata

references. This means that any evolution of the underlying API is taken into account by the generation engine and the program code remains unchanged.

Recently attributes are used in a number of enterprise frameworks. Attributes are normally employed to embed in the source code information that was previously specified in external configuration files or derived from conventions of source code elements. For example, in the EJB3 [10] specification, annotations on JavaBeans coexist with the legacy XML-based descriptors. Also, in JUnit version 4, instead of relying on the naming convention that the name of test-case methods must start with the string `test`, an annotation `@Test` is used. Annotations in these two frameworks (i) enhance the readability of the source code, in EJB3, the programmer must only look in a single file to get all the information for the EJB; and (ii) make the use of frameworks easier, in JUnit<sup>2</sup>, if the programmer misspells the `@Test` annotation, the Java compiler will flag the error, whereas this is not true for naming conventions. Annotations may also be used to directly represent DSLs. The aspect-oriented programming language AspectJ is a DSL closely based on Java, with the introduction of annotations in Java5, AspectJ migrated to being an annotation-based language by translating the DSL terms to equivalent annotations on pure Java code. By using AspectJ as a library rather than a language, the programmer can simplify her compilation process since both, aspects and base code, are processed by the same compiler.

The rest of the paper is organized as follows: in Section 2 we introduce Java 5 annotations as a platform for `@OP` and the Spoon annotation processor. In Section 3 we present our proposed framework for annotation validation, `AVal`; while in Section 4 we present two case studies: `Fraclet` (4.1) and an implementation of the JSR 181 (4.2). In Section 5 we evaluate the approach and in Section 6 compare it to related work. Finally conclude in Section 7.

## 2. Background

In this section we discuss Java 5 Annotations –with special attention to the validation capabilities provided within the language– and Spoon, our source code processor.

### 2.1. Java5 Annotations

In version 1.5, Sun included several language updates to Java, in particular, a metadata facility for program elements called *annotations*. Annotations in Java are a kind of typed metadata, that allows for `@OP`. They are defined in a way similar to interfaces, using the keyword `@interface`. Each annotation contains a number of **attributes** that can

<sup>2</sup><http://www.junit.org>

be primitive types (`int`, `float`, `String`), enumerations or classes represented as the return type of methods in the `@interface`. The semantics of the annotations can be implemented either at runtime (through the reflection API), or at compile-time (through the use of an annotation processor). An example of a Java5 annotation that allows programmers to specify which methods of a given class to test is shown in figure 1.

```
@Target({ElementType.CLASS,ElementType.METHOD})
public @interface Test{
    String name() default "";
}
```

Figure 1. Java annotation definition

Java5 annotations allow for a limited set of validations that restrict the source code elements on which the annotations can be placed. This is done by annotating the definition of the annotation (meta-annotating) with `@Target`. The `Target` annotation takes as argument an array with the elements on which the annotation can be placed. For example, the annotation defined in figure 1 can only be used on classes or methods. All other restriction on the usage of the annotations must be implemented by the `@OP` framework programmer. Depending on the way the `@OP` framework interprets the annotations, the misuse of an annotation will only be detected when the application is compiled or at runtime when the annotations are interpreted, if at all.

### 2.2. Spoon Annotation Processor

Spoon [11] is a source-code processor based on a meta-model of the program that models every code element, including statements and expressions. It relies heavily on generics to ensure type safe processing, and uses the concept of *processors* as units of program analysis and transformation. A processor is in essence an implementation of a visitor on the program’s model. In each visiting, the processor has complete (both read and write) access to the model. Special processors are `AnnotationProcessors` that declare the annotation in which they are interested, and the type of elements on which the annotation is applied.

## 3. AVal Annotation Validator

To provide a generic and extensible framework for the checking of annotated programs in Java, we have implemented an *Annotation Validator* (`AVal`) as a Spoon processor. `AVal` follows the idea that annotations should describe the way in which they should be validated, and that self validation is expressed by *meta-annotations* (`@Validators`). This idea of meta-annotations as a way to describe the rules

of use of domain level annotations is a generalization of Java's `@Target` annotation discussed in section 2.1.

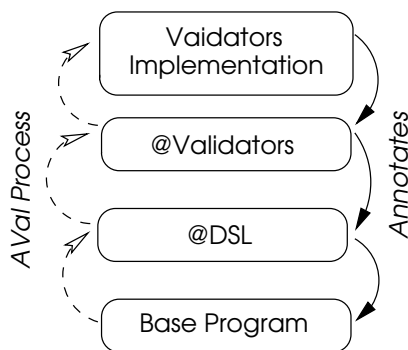
AVal's architecture is composed of four layers (figure 2):

**Base program** The (annotated) program that is to be validated. Elements of the program are annotated by attributes defined on the `@DSL` layer.

**Domain Specific (Annotation) Language (@DSL)** The domain specific annotations. Each annotation is meta-annotated by an AVal meta-annotation that expresses the rules for its validation.

**@Validators** Meta-annotations that encode the rules to validate annotations. Each `@Validator` represents a validation rule, and is itself annotated with the class that is responsible for the implementation of it.

**Validation implementation** A class implementing each validator. The class must implement the `Validator` interface, and it uses the Spoon compile-time model of base the program, `@DSL` annotation, and `@Validator` in order to perform the validation.



**Figure 2. AVal Architecture**

AVal is implemented as a Spoon source code preprocessor that is executed before the code generation or compilation phase in a `@OP` framework. It traverses the base code looking for domain specific annotations. Each time it finds an annotated element, it checks the model of the annotation's declaration to see if it has any `@Validators`. In case the annotation has one or more validators, the tool executes the each validator's implementation in the order in which they are defined. As a preliminary optimization, the validator's implementation is cached, so that if in the traversal of the program the same annotation is found twice, the correct validator implementation is executed without processing the annotation's definition again.

```

@Target({ElementType.METHOD})
@Inside(Documented.class)
public @interface WebLink{

    @URLValue String value();

}

@Implementation(URLValueValidator.class)
public @interface URLValue{}
  
```

**Figure 3. Example of the use of AVal @Validators: (@Inside and @URLValue)**

### 3.1. A small example

In order to better explain the nature of AVal validations we show a small example. Take a small `@DSL` composed of the annotations `@Document` and `@WebLink` that allows to document classes by means of external web pages. As validation rules we define:

- Only methods can be annotated with `@WebLink`
- Only methods within *Documented* classes (classes annotated `@Document`) can be annotated with `@WebLink`.
- The `@WebLink` annotation contains an `String` attribute that must be a valid URL.

In order to automatically check these rules, we annotate (figure 3) the definition of the `@WebLink` with two `@Validators`: `@Inside`, that checks that all `WebLinks` are inside code elements annotated `@Document`, and `@URLValue` that states that the value of the annotation should be a valid URL. The definition of these two `@Validators` will be explained in detail in the following sections.

When the AVal processor is executed, it visits the program's elements checking if they have a `@WebLink` annotation. When it finds one, AVal opens the definition of the `@WebLink` annotation and processes the `value()` element; it then creates an instance of the `URLValueValidator` class and executes the implementation of the `@Validator`. A detailed description of this process is found in section 3.3.

### 3.2. Generic Validators

As a starting point, AVal defines a number of *generic* validators that can be used regardless of the domain. These generic validators are divided into structural and value validators.

### 3.2.1 Structural Validators

Structural validators reason on the relationship between annotations, and, between an annotation and the code element it annotates.

**@AValTarget** This validator extends Java5's `@Target` meta-annotation allowing for finer control on what elements can an annotation be used. With it, it is possible to state that an annotation can only be placed on interfaces, for example. The `@Validator` defines a single attribute:

- `value()`: The type of the target element. This type is expressed by the class that represents the element in Spoon's compile-time model. For example `CtInterface.class` if it is an interface.

**@Inside** This validator states that the current annotation must occur within the lexical scope of another annotation. For example, a method annotation `@Foo` can only be used on methods that belong to classes that are annotated `@Bar`, then `@Foo` is *inside* `@Bar`. This `@Validator` defines a single attribute:

- `value()`: The type of the parent annotation (For example `Parent.class`)

**@Prohibits** This validator states that if a code element is annotated with the current annotation, then it prohibits the use of another (given) annotation. This `@Validator` defines a single attribute:

- `value()`: The type of the forbidden annotation.

**@Requires** This validator is the dual of the `@Prohibits`. It states that the current annotation must occur in elements that are also annotated with another annotation. This `@Validator` defines a single attribute:

- `value()`: The type of the required annotation.

### 3.2.2 Value Validators

Value Validators reason on the attributes of the annotations. They allow to overcome Java5 restriction on the types allowed for annotation's attributes (primitive types, Classes, etc). Also they allow for domain-specific semantic checks.

**@RefersTo** This validator states that the value of the attribute must refer (be equal) to the value of another annotation somewhere in the program. This `@Validator` defines the following attributes:

- `value()`: The target annotation type to which the current attribute refers to.

- `attribute()`: The attribute in the target annotation type to which the current attribute refers to.

**@Matches** This validator applies on attributes of type `String`, and it checks that the values of the attribute match the provided Java regular expression. This `@Validator` defines a single attribute:

- `value()`: The regular expression.

**@Unique** This validator checks that a given attribute value is unique for a given program. That is, two program elements annotated with the same annotation type cannot have the same value in an `Unique` attribute. This `@Validator` defines no attributes of its own.

```
@A("http://localhost/")
public class Foo{
//...
}
```

a. Base Program

```
public @interface A{
    @URLValue
    String value();
}
```

b. @DSL annotation

```
@Implementation(URLValueValidator.class)
public @interface URLValue{}
```

c. @Validator

```
public class URLValueValidator
    implements Validator<URLValue>{

    public void check(
        ValidationPoint<URLValue> vp){

        String attribName = vp.getDslElementName();
        String value =
            (String)vp.getDslAnnotationValue(attribName);

        try{
            new URL(value);
        }catch(MalformedURLException ex) {
            //Report error
        }
    }
}
```

d. Validation implementation

Figure 4. Custom Validator

### 3.3. Extending Validators

We have found, as it will be discussed in section 4, that the generic validators shown before cover many validation needs. However, there are cases in which it is difficult or even impossible to translate a given domain rule into generic validators; for these cases we have implemented a way to extend the `@Validator` set for a particular domain.

New validators require two things: a new `@Validator` annotation, and its corresponding implementation. `@Validators` are normal Java annotations that are themselves annotated with their corresponding implementation. The implementation of a `@Validator` is a class that implements the `Validator` interface parametrized by the type of the `@Validator`. This interface defines a `check` method that is up called whenever the validated annotation is found. Validator implementations have access to the complete meta-model of the program, in particular to the annotated base program, the `@DSL` annotation and annotation definition, and the `@Validator`. These elements of the meta-model are encapsulated in a `ValidationPoint` object.

To illustrate the process of definition of a new validator, suppose that a domain rule requires that the value of the attribute `value()` of the annotation `A` is a valid URL (fig. 4b.). To this end, we define a `@Validator URLValue` (fig. 4c.). The class `URLValueValidator` is responsible for checking that the values of the `@A` are in fact URLs (fig. 4d.).

By providing access to the program's compile-time model and using Java, it is possible to implement complex validators, for example `@RefersTo`, that take into account the complete program, or validate with regard to external resources, such as checking an XML specification, or database schema.

## 4. Case Studies

In this section we present two Attribute-oriented specifications using Java5 annotations. Fraclet [12], an annotation framework for the Fractal component model [1] and the JSR181 for definition of web services [16]. We show how our proposal allows for a clear definition of the rules of use of these specifications.

### 4.1. Fraclet

Fraclet is an annotation framework for the Fractal component model. The Fractal component model defines the notions of *component*, *component interface*, and *binding* between components. Each of these main notions is reflected in the `@DSL` defined by Fraclet. There are two implementations of Fraclet, one using XDoclet2, and the other

one using Java5 annotations and Spoon annotation processor. The annotations defined by Fraclet/Spoon are summarized in table 1.

The rules for the use of each of the annotations in Fraclet/Spoon are as follows:

**@FractalComponent** A Fractal component in Fraclet/Spoon is a Java **class** that defines a number of component attributes, bindings and operations. The `@Target` annotation provided by Java only allows to define that the annotation can be placed on types (classes or interfaces), therefore, the `@Validator @AValTarget` is used to restrict the Fractal components to being only classes. The complete definition of the annotation is shown below.

```
@AValTarget (CtClass.class)
public @interface FractalComponent {
    String controllerDesc() default "";
}
```

**@FractalItf** A Fractal business interface is a Java interface that defines a set of related operations in a component. The interface must contain a name that is unique for the application, and it must define if the interface is optional, and its cardinality. `@Validators` are provided to check all these rules:

```
@AValTarget (CtInterface.class)
public @interface FractalItf {
    @Unique String name();
    Class signature() default None.class;
    @Matches("(singleton|collection)")
    String cardinality() default "singleton";
    @Matches("(mandatory|optional)")
    String contingency() default "mandatory";
}
```

**@FractalAC** A field annotated as `FractalAC` describes an attribute of the Fractal component, therefore, only fields that belong to a Fractal component class are allowed to be annotated `@FractalAC`. Also, since Fractal attributes and Fractal bindings are both represented using fields, it makes no sense to annotate a single field with both `@FractalAC` and `@FractalBC`. `@Validators` for these rules are included in the definition of the annotation:

```
@Inside (FractalComponent.class)
@Prohibits (FractalBC.class)
@Target (ElementType.FIELD)
public @interface FractalAC {
    String argument() default "";
    String value() default "";
}
```

**@FractalBC** A Fractal binding represents a binding between a component and a Fractal interface. The binding is represented as a field in a Fractal component class, and

Annotation	Location	Parameter	Description
@FractalComponent	Class	<i>controllerDesc</i>	Annotation to describe a Fractal component.
@FractalItf	Interface	<i>name, signature, cardinality, contingency</i>	Annotation to describe a Fractal business interface.
@FractalAC	Field	<i>argument, value</i>	Annotation to describe an attribute of a Fractal component.
@FractalBC	Field	<i>name, signature, cardinality, contingency</i>	Annotation to describe a binding of a Fractal component.
@FractalImportedInterface	Class	<i>interfaces</i>	Annotation to specify that the component provides a server interface which is not annotated with a @FractalItf.
@FractalIRC	Field	-	Annotation to get the component part interface

**Table 1. Overview of Fraclet annotations**

therefore, is only valid in fields of classes annotated with @FractalComponent. It defines the name of the Fractal interface that is bound to (which must exist in the program), as well as the signature, cardinality, and contingency of the binding. These last three attributes follow the same rules than those of @FractalItf.

```
@Inside(FractalComponent.class)
@Prohibits(FractalAC.class)
@Target(ElementType.FIELD)
public @interface FractalBC {
    @RefersTo(value = FractalItf.class, attribute="name")
    String name();

    Class signature() default None.class;

    @Matches("(singleton|collection)")
    String cardinality() default "singleton";

    @Matches("(mandatory|optional)")
    String contingency() default "mandatory";
}
```

**@FractalImportedInterface** Fractal components implement interfaces that may not be Fractal business interfaces, but that still need to be exposed in the component; for example `java.lang.Runnable`. These interfaces are declared as *imported interfaces* in the definition of the Fractal component, therefore, it makes no sense to annotate a class with @FractalImportedInterface if it is not a Fractal component. Note that the `interfaces()` attribute is an array of @FractalItf, and therefore it is checked using the rules defined for Fractal business interfaces.

```
@Requires(FractalComponent.class)
public @interface FractalImportedInterface {
    FractalItf[] interfaces();
}
```

## 4.2. JSR 181

The JSR181 [16] is a specification for the description of web services using pure Java objects. The JSR defines a set of annotations and their mapping to the XML-Based Web Service Description Language. In section 2.5.1 of the specification, it is stated that implementations of the JSR must provide a validation mechanism that performs the semantic checks on the Java Bean web service definition. Table 2 summarizes the six of the ten annotations defined by the JSR.

Rules defined for the JSR describe restrictions not only on the use of the annotations, but also on certain properties of the annotated elements, for example that the web service implementation must not define a `finalize()` method, or that a *one-way* operation must have no return value. For this domain specific restrictions we extend the validation framework with a new @Validator for each annotation. This @Validator encapsulates all checks regarding the contents of the annotated element. The selected annotations of the @DSL are discussed below.

**@WebService** This annotation marks a Java class as a service implementation bean, or a Java interface as an endpoint interface. As the same annotation is used to describe two entities: service implementation and endpoint interface, the constraints on the annotated element vary depending on if the annotation is placed on a Java class or an interface. Regardless of where the annotation is placed, the `wSDLLocation()` attribute must be a valid URL.

If a class is annotated @WebService, it must be an outer class and it must not be `final` nor `abstract`, it must also define a default public constructor. These rules are validated by the `ValidWebServiceBean` @Validator. If an interface is annotated @WebService, it is required that the interface is `public` and the annotation is not allowed to define values for the `serviceName()` and `endPointInterface`. These rules are validated by the



Annotation	Location	Parameter	Description
@WebService	Class, Interface	<i>name, targetNamespace, serviceName, wsdlLocation, endpointInterface</i>	Class or Interface defining a web service
@WebMethod	Method	<i>operationName, action</i>	Method exposed as a web service operation
@OneWay	Method	–	Indicates that a given web server operation has only input messages and no output.
@WebParam	Method Parameter	<i>name, targetNamespace, mode, header</i>	Maps an individual operation parameter to a web service message
@WebResult	Method	<i>name, targetNamespace</i>	Maps the operation's return value to a web service result
@HandlerChain	Class, Interface	<i>file, name</i>	Associates an externally defined handler chain to a web service

**Table 2. Overview of JSR-181 annotations**

ValidEndPointInterface @Validator.

```
@Target( { ElementType.TYPE })
@ValidWebServiceBean
@ValidEndPointInterface
public @interface WebService {
    String name() default "";
    String targetNamespace() default "";
    String serviceName() default "";
    @URLValue
    String wsdlLocation() default "";
    String endpointInterface() default "";
};
```

**@WebMethod** This annotation marks a method as being a web operation for the web service. The method must be public, and its parameters and return type conform to the rules defined in the JAX-RPC specification [3]. The checks of the signature of the method are implemented in the ValidWebOperation @Validator.

```
@Target( { ElementType.METHOD })
@Inside(WebService.class)
@ValidWebOperation
public @interface WebMethod {
    String operationName() default "";
    String action() default "";
};
```

**@Oneway** This annotation states that a given WebMethod has only an input message, and no return value. The methods annotated @Oneway cannot declare checked exceptions, or define IN or INOUT parameters. The checks on the signature of the web methods are carried out by the ValidOneway @Validator.

```
@Requires(WebMethod.class)
@ValidOneway
public @interface Oneway {
};
```

**@WebParam** This annotation defines the properties for parameters of web methods. The specification does not define particular rules about this annotation other than that it must be defined only on parameters of web methods.

```
@Target( { ElementType.PARAMETER })
@Inside(WebMethod.class)
public @interface WebParam {
    public enum Mode {
        IN, OUT, INOUT
    };

    String name() default "";
    String targetNamespace() default "";
    Mode mode() default Mode.IN;
    boolean header() default false;
};
```

## 5. Evaluation

As discussed in previous sections, AVal has been applied to two @OP frameworks. In the case of Fraclet/Spoon we were able to check all the rules using only generic validators (3.2), without having to implement Fraclet-specific ones. Given that Fraclet/Spoon uses the same source-code transformation tool than AVal, we were able to integrate the @Validators seamlessly to the @OP framework. AVal is included in the latest release of Fraclet/Spoon.

In the case of the JSR181, rules about the relationship between annotations and rules that restrict the values of the annotation's attributes are encoded using generic validators, while restrictions on the program elements are validated by custom-made @Validators. It is interesting to note that the restrictions on the program elements are not directly related to the @DSL itself, but to the domain in which the annotations are used. Indeed, the rule that states that a class that implements a web service must not be final is independent of the way in which the class is marked as a web service. In this regard, the @Validators not only check annotations, but also domain related restrictions.

A weakness identified during the case studies is that, in order to attach the `@Validators` to the annotations defined in the `@DSL`, the developer needs to modify the definition of the annotations. This presents a problem when the `@DSL` is part of an external library (such as the Reference Implementation of the JSR181). This fact restricts the use of `AVal` to cases in which the `@OP` framework developer is also in charge of the `@DSL` definition, we do not believe this to be a too strong assumption.

## 5.1. Future Work

For the continuation of `AVal`, we expect to apply it to more complex `@OP` frameworks such as EJB3 (which defines more than fifty persistence annotations) and AspectJ5. The *AValidation* of these frameworks will allow us to verify and expand the number of generic validators, as well as to test the performance of the approach against large applications.

We would like also to explore the possibility of extending `AVal` to non-annotation based frameworks. The idea would be to annotate elements in the framework with `@Validators` that would check that they are correctly used. The tests would be similar as those performed by the `@ValidEndPointInterface` and `@ValidWebServiceBean` defined in Section 4.2.

## 6. Related Work

**Static Validation** Static validators allow developers to check properties of their code that go beyond of that what is provided by normal compilers. Lint [9] is one of the first tools to provide such checks by relying on (lightweight) static analysis. To reduce the amount of noise (false positives) that is normally generated by Lint-like tools, LCLint [5], and later Splint [6], guide the validation of programs through annotations (stylized code comments) that explicit programmer assumptions and intents. This use of annotations is comparable that of `AVal`'s `@Validators`.

In [7], Hedin proposes an extensible, attribute-based static validator. In it, the grammar of a language is extended to check that custom programming conventions are followed. These extensions are similar in spirit to those possible with `AVal`; nevertheless, they lack the modularity and cohesion of implementing each extension in a separate class as is done in `AVal`, since the extension of the grammar is done by attributing each individual node of the AST and then acting upon these attributes, thus lacking locality.

By regarding validation as a crosscutting concern in a program's code, it is possible to encode it by means of Aspect Oriented techniques, this has been explored by Shomrat et. al. in [13]. Nevertheless, in an Aspect Oriented language such as AspectJ[8], no extra reflection facilities

are provided, so the validation programmer must rely only on Java reflection which does not reify the body of methods and, since reflection is implemented at runtime, the `@OP` framework must be modified so that annotations are kept until runtime (using a special Java meta annotation). This restricts the domain of validations that can be performed

**Annotation Validation** Previous to the introduction of annotations in Java, XDoclet [15] relied on modified javadoc comments called tags to specify metadata for program elements. In XDoclet2, a form of tag validation is performed by tagging the tag definitions. The set of validations is fixed, and no special facilities are provided for extending them.

Cepa and Mezini's work [2] follows an approach similar to ours. They propose a meta annotation for the custom attribute facility in the .Net framework. However, they concentrate on *dependencies* between annotations (what we call structural validations 3.2.1) and do not foresee extensions to their model. In a later work [4], they propose an approach that is more general since it allows to validate constraints between different artifacts in the system (source code, configuration files, etc.). However, these constraints are expressed by means of a separate XML-based language, which in our opinion, goes against the idea of `@OP` which strives to reduce the use of external configuration files as much as possible.

## 7. Conclusion

We have presented `AVal`, an approach based on meta-annotations for the validation of the use of annotations in `@OP` Java applications. `AVal` is a declarative, expressive and extensible way to define and reuse validations of annotations. Also, it enhances the readability of `@DSL` source code definitions by embedding semantic information in their declaration. We have provided as case studies two `@OP` frameworks: `Fraclet` and the JSR181 for web services, and shown how to use `AVal` to include syntactic as well as semantic checks in them.

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