

# A Value Analysis for C Programs

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CEA List

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mp2 []]] = -(1 << [NB - 1]] set if the function of the second part of



# Frama-C: a static analysis framework for C

#### A static analyzer is a plug-in





- A static analyzer is a plug-in
- Frama-C and plug-ins are written in OCaml
- Plug-ins can be distributed independently from Frama-C core
- No recompilation of core to integrate a new plug-in
- Plug-ins provide services to each other



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Think "The Gimp" or "Eclipse" http://frama-c.cea.fr/



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# Initial focus of Frama-C: verification of critical embedded code. Existing plug-ins are **correct**



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#### Establish safety properties

if codebase in the right C subset,  $\ldots$ 



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Implementing a heuristic method as a Frama-C plug-in?



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Sure, go ahead...



mp2[]]] = 1 << NME = 10; eVer (f mp1[]]) = (1 << NME = 11) (mp2[]]] = 11 << NME = 11 << NME = 11 = (1 ke mp2[]]] = mp2[]]] = mp2[]] = mp2[] = mp2[] = mp2[] = mp2[] = mp2[]] = mp2[] = mp2[] = mp2[] = mp2[] = mp2[] = mp2[]] = mp2[] = mp2



### Value Analysis plug-in

Abstract Interpretation (comparable to Astrée or PolySpace) Emphasis on pointers, pointer casts, handling of statically allocated chained structures



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Results recorded for each statement for use by other plug-ins. 100 000 lines of embedded code — 10GiB of memory (peak)



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### **Clarification Questions**





**Controversial Statement** 

Attention C compilers authors:

Nobody except you cares about your speed benchmarks.



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Attention C compilers authors:

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Current trend in the C compilation world: take advantage of any little bit of undefinedness in the standard.

- signed arithmetic overflows not giving 2-complement results
- char\* assumed not to alias with an int\*

tmp2[]]]=r[<<</Net |Net[Net[Net]]] =r[<</Net|Net[Net]] (mp2]]] =r[<</Net|Net[Net]] =r[<</Net|Net[Net]] =r[Net][]] =r[Net][]] =r[Net][] (mp2]]] (mp2]] (mp2]]







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