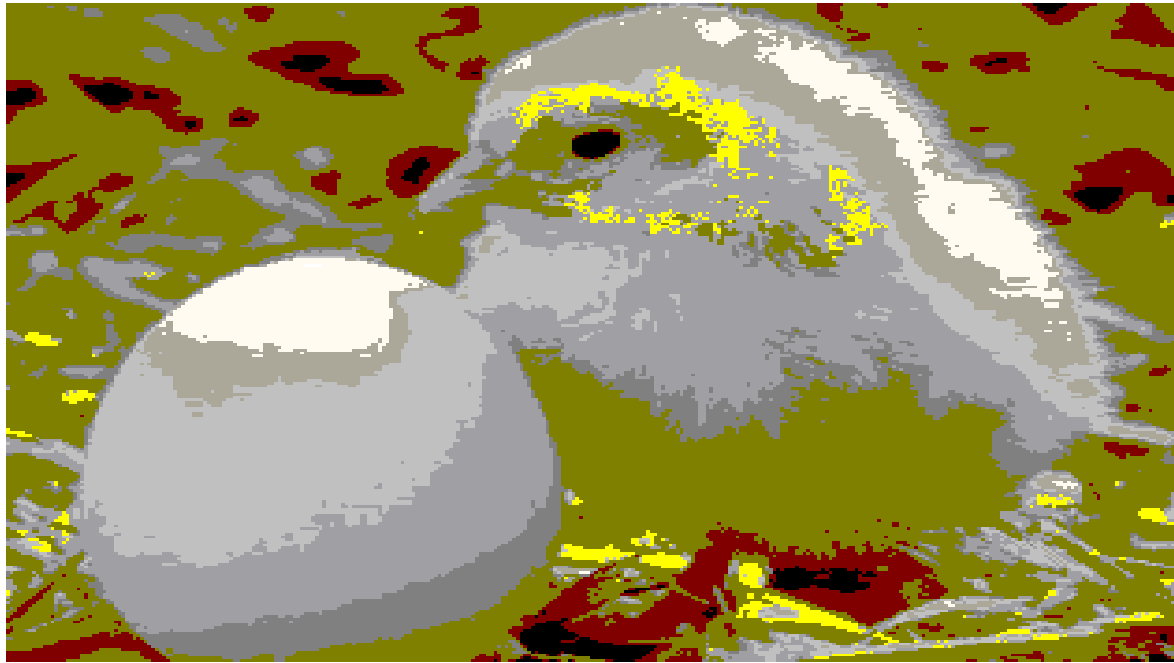


# Properties of Slicing Definitions

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Is  $S$  a slice of  $P$



# A Simplistic Definition

- Program **S** is a slice of Program **P** taken with respect to Variable **V** if when run on initial state  $\sigma$ 
  1. **P** terminates
  2.  $\mathcal{M} \mathbf{P} \sigma \upharpoonright_{\mathbf{V}} = \mathcal{M} \mathbf{S} \sigma \upharpoonright_{\mathbf{V}}$   
(**P** and **S** compute the same final value for **V**)

# An Example

$x := 1$

*is a slice of*

$x := 1$

$y := 2$

# And when **P** fails to terminate?

is

*x := 1*

a slice of

*x := 1*

*while true do skip od*

or

*while true do skip od*

*x := 1*

# Goal: Characterize Slicing using *Properties*

Two kinds

- Floor
- Ceiling

# Floor Requirement

- A “floor” requirement describes  
*“must includes”*
- For example  
a program must be a valid slice of itself

# Ceiling Requirement

- A “ceiling” requirement describes  
*“must avoids”*
- For example  
a non-terminating program cannot be  
a slice of a terminating program



# More formally

- Define a slicing relation as a set of triples  $\langle \mathbf{S}, \mathbf{P}, \mathbf{X} \rangle$ 
  - $\mathbf{S}$  is a slice of  $\mathbf{P}$  taken with respect to  $\mathbf{X}$
- Slicing relation  $\mathbf{R}$  satisfies
  - floor requirement  $\mathbf{F}$  iff  $\mathbf{F} \subseteq \mathbf{R}$
  - ceiling requirement  $\mathbf{C}$  iff  $\mathbf{R} \subseteq \mathbf{U} - \mathbf{C}$

# A Floor Example

## *Weaken Criterion*

If  $\langle \mathbf{S}, \mathbf{P}, \mathbf{X} \rangle \in \mathbf{R}$  and  $\mathbf{X}' \subset \mathbf{X}$   
then  $\langle \mathbf{S}, \mathbf{P}, \mathbf{X}' \rangle \in \mathbf{R}$

## Example

“if  $\mathbf{S}$  preserves the final value of  $\mathbf{a}$  and  $\mathbf{b}$   
then it preserves the final value of  $\mathbf{a}$ ”

$\langle \mathbf{S}, \mathbf{P}, \{\mathbf{a}, \mathbf{b}\} \rangle \in \mathbf{R}$  requires  $\langle \mathbf{S}, \mathbf{P}, \{\mathbf{a}\} \rangle \in \mathbf{R}$

# Another Floor Example

*Truncation*

$\langle S1, S1;S2, \{X\} \rangle$  must be included  
if  $S2$  is *X-preserving*

Statement  $S$  is *X-preserving* if it does not  
change the value of  $X$   
(bit of a simplification)

# Final Floor Example

## Ditchability

The slicing relation should allow deletion of any code that does not affect the value of any variable of interest

# Ceiling Examples

## *Termination Preserving*

Must avoid  $\langle \mathbf{S}, \mathbf{P}, \mathbf{X} \rangle$  when  $\mathbf{P}$  terminates and  $\mathbf{S}$  does not

low, all slices of  $\mathbf{P}$  must terminate when  $\mathbf{P}$  terminates

# Ceiling Examples

## *Non-termination Preserving*

Must avoid  $\langle \mathbf{S}, \mathbf{P}, \mathbf{X} \rangle$  when  $\mathbf{P}$  diverges  
and  $\mathbf{S}$  does not

low, all slices of  $\mathbf{P}$  must diverge when  
 $\mathbf{P}$  diverges

# Ceiling Examples

## *Behaviour Preserving*

$\langle \mathbf{S}, \mathbf{P}, \mathbf{X} \rangle$  requires that when  $\mathbf{P}$  terminates  
 $\mathbf{S}$  also terminates

$\forall \mathbf{x} \in \mathbf{X}$   $\mathbf{x}$ 's final value is the same for  $\mathbf{S}$  and  $\mathbf{P}$

# Properties allow Characterizing Slicing Approaches

- **Weiser's Definition**
  - Slice **S** of Program **P** is an executable program obtained from **P** by removing statements, such that **S** replicates part of the behaviour of **P**
- Definition is Termination Preserving
- Definition is **not** Non-termination Preserving



# Slicing Defined using a Lazy Semantics

- A lazy or demand semantics can be used to define the semantic of a slice

$x = 1$

$y = 2$

$z = x * 3$  -- demands  $x$  but not  $y$

- Does not satisfy *Behaviour Preservation* nor *Truncation*

# Slicing Defined using Semi-refinement

For Programs  $S1$  and  $S2$

$S2$  is a *semi-refinement* of  $S1$  provided

1. If  $S1$  terminates on state  $\sigma$ , then  $S2$  also terminates on  $\sigma$  with the same set of possible final states
2. If  $S1$  does not terminate, then  $S2$  can do *anything at all*

# Slicing Defined using Semi-refinement

Using semi-refinement satisfies

- Weaken
- Identity
- Termination Preservation
- Ditchability
- Truncation
- but **not** Non-termination Preserving

# Properties bound possible slicing relations

For example, the two properties

*Behaviour Preservation*

and

*Truncation*

are sufficient to restrict possible slicing  
relations to just one: semi-refinement

# Property Relations

Non-termination preservation is incompatible with ditchability.

# Controversial?

Let **P** be

$a = 42$

$a = 4$

while true do skip

Slice **P** at the end with respect to  $a$   
**S1:**

$a = 42$

is a slice of **P**

as is **S2:**

$a = 46$

*(a poll of attendees showed little controversy as none thought that **S1** nor **S2** was a slice of **P**)*